CAB301 Assignment 2

Empirical Comparison of Two Algorithms

Student name: John Santias, Shivam Sachdeva

Student no. n9983244, n9642587

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**Summary**

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**1 Description of the Algorithm**

**1.1 Brute Force Median Algorithm**

The brute force median algorithm returns the median value in a given array. Typically, in a sorted list the median would always be in the middle, but by having an unordered list it can be harder to find the value. “Brute force” means the algorithm will try to find many possible ways to find the median value of a given array. This algorithm’s pseudocode can be found in **Appendix A**. Referring to this pseudocode, firstly, the algorithm assigns the value in the middle of the list to a variable named ‘*k*’. This is assuming that we do have the correct median value seeing as the list is ordered or not. When selecting a median inside an uneven array size (decimal rounded up), the program will choose the value in the middle, but with an uneven sized array, it will choose the left of the two middle elements (). No matter the selected median is correct or not, the following for-loop runs to double check it. This for-loop works its way through the array from left to right. Each array value is used as a pivot to be compared with the other values. This loop starts off with assigning the two variables, ‘numsmaller’ and ‘numequal’, with a value of zero. Afterwards, another for-loop is executed to use each array value for comparison with the pivot. Inside this inner loop, there are two comparisons. The first comparison determines whether there is a value in the array less than the current pivot. If it is, then the variable ‘numsmaller’ increments by one. Otherwise, the second comparison determines if the values are equal then it increments the variable ‘numequal’ by one. Breaking back out to the original for-loop, a final comparison determines if it has chosen the correct median value by seeing that the value of ‘k’ is greater than the value of ’numsmaller’ and it is less than and equal to the total value of ‘numsmaller’ and ‘numequal’. In other words, it determines if the value of ‘k’ is the median value of all values in the array, seeing that there are even numbers on the left and the right side of the chosen median or there is one more value on the right compared to the left side of the chosen median. If this comparison is not met, the for-loop increases by one and executes another iteration of the whole process.

**1.2 Median Algorithm**

The Median algorithm is used to find the median of an array of elements. There are three separate procedures in this algorithm. This algorithm is special case of Johnsonbaugh and Schaefer’s version of the selection problem algorithm. The pseudocode of the algorithm is shown in **figure 2 of Appendix A**. The main procedure called median takes the array as a parameter. If the array contains only one element, the algorithm returns that element as the median whereas if the array contains more than one element, the procedure calls the second procedure called select. The select procedure takes four parameters: array, first index value of array, index of median element and the index of last element of array. This procedure calls the partition procedure which takes the first element of array as a pivot and compare it with all the other values of array. Every time the pivot value is greater than the other element, the pivotloc variable increases by one and swaps the element around the pivot. After the for-loop is executed, the pivot is placed on position where it should be if the array was sorted. This procedure returns the pivotloc variable. The value returned from partition procedure is stored in pos variable of Select procedure and compared with the index of median element which is equal to n/2 where n is number of elements in array. If the value matches, the select procedure returns the median of array and in case value does not matches, the select procedure is called again but not with the subarray. Now again partition procedure is called to do partitioning on subarray. The returned value of partition procedure is compared again and this process goes on until the value matches with the index of median element. Basically, Select algorithm is a recursive method which is called many times until the median is found. Every recursive call involves an array slice whose size is reduced every time.

**2 Theoretical Analysis of the Algorithm**

**2.1 Choice of Basic Operation**

The basic operation of Brute force median algorithm is the if and else-if conditions within nested for loops. Every array element is compared to itself and to the other elements. The block of if and else if condition checks whether the element is greater than or equal to other array elements. If the array element is greater than the other element, the ‘numsmaller’ variable increments where if two elements are equal, the ‘numequal’ variable increments. This block of if and else if condition is considered as the one basic operation which determines the efficiency of brute force median algorithm. The basic operation of the Median algorithm is the if condition within the for-loop of the partition procedure which checks whether the pivot is greater than the other elements. Basically, how many times pivot is checked against other elements or in other words how many times for-loop executes for an array determines the efficiency of the Median algorithm. The partition procedure can be called recursively many times until the median is found. The basic operation chosen for both the algorithms performs same function which is comparing element with the other array elements.

**2.2 Choice of Problem Size**

The problem size given to our algorithms ranges from one to a hundred. Incrementing the problem size by one. We implemented a for-loop re-iterating 100 times to generate arrays with sizes from one to a hundred, then passing the array to the two median algorithms to solve. Each element value in the Array has a random generated number between zero and ten thousand. This is because we are dealing with average-case efficiency of our algorithms, we don’t want our array elements to be in order or in descending order. Both of these are best and worst-case efficiencies. As we are comparing two algorithms, we will be giving the algorithms the same input sizes and same element values.

**2.3 Average-case efficiency**

The for-loop of the Brute Force Algorithm iterates from to

… The outermost for-loop of brute force algorithm iterates from i = 0 to i =  n-1 and the inner for-loop iterates from j = 0 to j = n-1. The basic operation is within inner for-loop.

Mathematically,

Why does Brute force have a quadratic efficiency?

**3 Methodology, Tools, and Techniques**

**3.1 Programming Environment**

1. We decided to implement the algorithms and perform the experiments in the C# programming language. As having used it before, we found it easy to run experiments and retrieve results.
2. The experiments were performed on a Windows 10 PC. We used the software called visual studio community, to run our C# experiments. We used C#’s random number generator to create random values inside our array and its system timer to record the program’s execution time.
3. A Microsoft Excel spreadsheet was used for recording our results and producing graphs. This software will help organise our experiment results in separate columns and rows. It will also help us calculate the average result of basic operations and time efficiency for each size input automatically. Using the results, we were able to generate a graph using Excel’s graph function.

**3.3 Implementation of the Algorithms**

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**3.4 Generating Test Data and Running the Experiments**

**4 Experimental Results**

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**4.1 Functional Testing**

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**4.2 Average-Case Number of Basic Operations**

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**4.3 Average-Case Execution Time**

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**References**

[1] Levitin, A., 2011. *Introduction to the Design and Analysis of Algorithms*. Addison-Wesley.

**Appendix**

Figure 1



Figure 2